

Optimization of hole-making operations for injection mould using particle swarm optimization algorithm

A. M. Dalavi^{a*}, P. J. Pawar^b and T. P. Singh^a

^aDepartment of Mechanical Engineering, Symbiosis Institute of Technology, Symbiosis International University, Gram Lavale, Mulshi, Pune, India 412115

^bDepartment of Production Engineering, K. K. Wagh Institute of Engineering Education and Research, Nashik, India

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ABSTRACT

Optimization of hole-making operations plays a crucial role in which tool travel and tool switch scheduling are the two major issues. Industrial applications such as moulds, dies, engine block etc. consist of large number of holes having different diameters, depths and surface finish. This results into to a large number of machining operations like drilling, reaming or tapping to achieve the final size of individual hole. Optimal sequence of operations and associated cutting speeds, which reduce the overall processing cost of these hole-making operations are essential to reach desirable products. In order to achieve this, an attempt is made by developing an effective methodology. An example of the injection mould is considered to demonstrate the proposed approach. The optimization of this example is carried out using recently developed particle swarm optimization (PSO) algorithm. The results obtained using PSO are compared with those obtained using tabu search method. It is observed that results obtained using PSO are slightly better than those obtained using tabu search method.

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1. Introduction

In machining process of many industrial parts such as dies and moulds, operations like drilling, reaming or tapping account for a large segment of process. Generally, a part, for e.g. a plastic injection mould may have many holes with different diameters, surface finish, and maybe various depths. If the diameter of hole is relatively large, a pilot hole may have to be drilled first using a tool of smaller diameter and then enlarge it to its final size with a larger tool, which could be followed by reaming or tapping whenever essential. For hole H_3 , as shown in Fig. 1, there could be four different combinations of tools:(A,B,C), (A,C), (B,C), and (C). The selection of tool combinations for each hole directly influences on the optimum cutting speeds, required number of tools switches, and tool travel distance (Kolahan & Liang, 2000).

* Corresponding author. Tel: +91-9921517645
E-mail: amol.dalavi83@gmail.com (A. M. Dalavi)

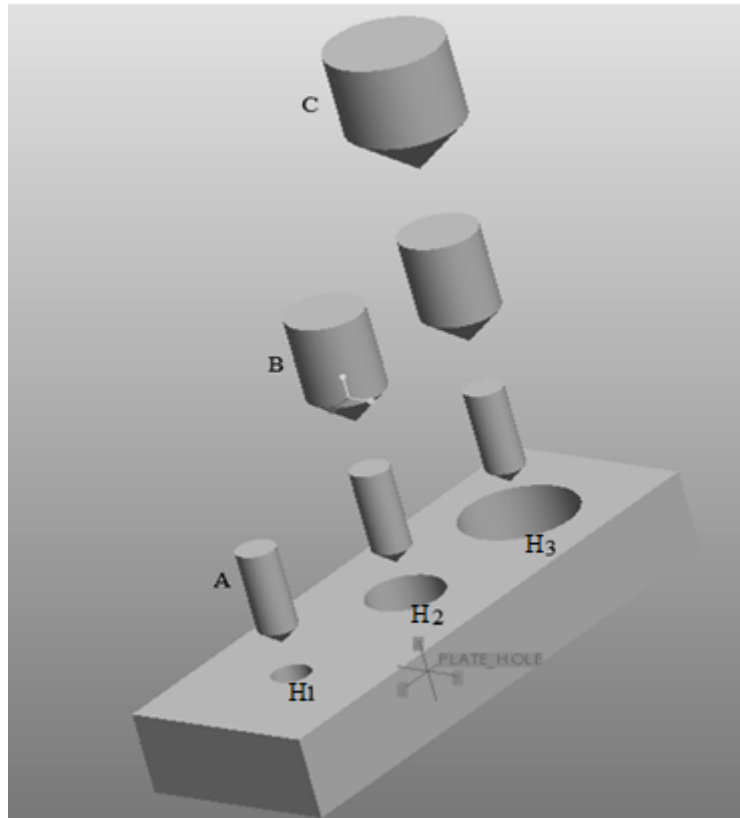


Fig. 1. Image showing various tool combinations required to drill a hole on workpiece

Tool switch and tool travel from one position to another takes more amount of machining time in machining processes. Usually 70% of the overall time in machining processes is spent on movements of tools and part (Merchant, 1985). To reduce the tool travel, the spindle is not moved until a hole is completely drilled using several tools in different diameters, which increases tool switching cost. On the other hand, to reduce tool switching cost, it may be used to drill all possible holes which, in turn, increases the tool travel cost. Luong and Spedding (1995) addressed the process planning and cost estimation of hole-making operations by developing a generic knowledge based procedure. Castelino et al. (2000) reported an algorithm for minimizing airtime for milling by optimally connecting various tool path segments. In their work, a problem was formulated as a generalized travelling salesmen problem and it was solved using a heuristic method. Kolahan and Liang (2000) introduced a tabu search approach to reduce the overall processing cost of hole-making operations. Alam et al. (2003) presented a practical application of computer-aided process planning (CAPP) system to reduce the overall processing time of injection moulds. Genetic algorithm (GA) was used for optimizing the selection of machine tools, cutting tools, and cutting conditions for different processes. Abu Qudeiri et al. (2007) used genetic algorithm to find the optimal sequence of operations which gives the shortest cutting tool travel path (CTTP).

Jiang et al. (2007) reported a stochastic convergence analysis of the parameters $\{\omega, C_1, C_2\}$ of standard particle swarm optimization (PSO) algorithm. Shi et al. (2007) presented a novel PSO based algorithm for solving the travelling salesman problem (TSP). They compared their proposed algorithm with existing algorithms and found that PSO could be used for solving large size problems. Zhang et al. (2008) presented an improved PSO algorithm (IPSO) based on the “all different” constraint to solve the flow shop scheduling problem with the aim of minimizing make span. Guo et al. (2009) developed a problem on integration of process planning, scheduling of manufacturing field using PSO algorithm. Shao et al. (2009) used a modified genetic algorithm based approach to integrate the process planning and scheduling of manufacturing systems in order to achieve an improved performance. Zhang and Zhu (2011) proposed two models of PSO algorithm; one is based on value exchange and the other based on

order exchange. Chandramouli et al. (2012) reported sheep flock heredity algorithm (SFHA) and artificial immune system (AIS) for reducing time of the scheduling of machines, an automated guided vehicle (AGV) and two robots in a flexible manufacturing system. Bhongade and Khodke (2012) proposed two heuristics for solving assembly flow shop scheduling problem. Shahsavari Pour et al. (2013) presented genetic algorithm for solving the flow shop scheduling problem.

Ghaiebi and Solimanpur (2007) applied ant colony optimization (ACO) algorithm for optimizing the sequence of hole-making operations of industrial part. Hsieh et al. (2011) used immune based evolutionary approach to find the optimal sequence of hole-making operations. Tamjidy et al. (2014) presented an evolutionary algorithm to reduce the tool travel and tool switching time during hole-making operations based on geographic distribution of biological organism.

It is revealed from the literature that non-traditional methods such as tabu search (TS), genetic algorithm, ant colony algorithm, immune algorithm (IA) etc. were used to solve the problem of optimization of hole-making operations. However, pure tabu search that uses only one solution can easily neglect some promising areas of the search space, and may also not find optimal or exact solution. Most commonly used advanced optimization techniques are the implementation for genetic algorithm in manufacturing optimization. Genetic algorithm (GA) gives near optimal solution for complex problems (Rao, 2011) and it requires more parameters (Elbeltagi et al., 2005). In ACO algorithm, convergence is slow due to pheromone evaporation and it tends to use more CPU time (Elbeltagi et al., 2005). Immune based evolutionary approach requires more parameters. Hence it is necessary to use non-traditional optimization algorithm, which gives correct solution for complex problems (Rao, 2011). From literature it is found that recently developed optimization algorithm known as PSO could be used due to its simplicity, easy implementation and high convergence rate (Coello et al., 2011). In this work an attempt has been made by using PSO to reduce overall processing cost of hole-making operations through determination optimal sequence for hole-making operations.

2. Formulation of an optimization model

In order to reduce the overall processing cost of hole-making operation, the following optimization model is formulated based on the analysis given by Kolahan and Liang (2000) considering the following components of overall processing cost:

2.1. Tool travel cost

Tool travel cost is the cost of moving the tool from its previous location to the current drilling location.

2.2. Tool switch cost

It occurs whenever a various tool is used for the next operation. If the required tool type is not available on the spindle for machining operation, then the required tool must be loaded on the spindle prior to performing a machining operation.

2.3. Tool and machining costs

Tool cost includes the new tool cost and the cost of machine down time required to replace the tool. Machining cost comprises the operating cost and the machine overhead cost. The combined tooling and machining costs when tool type m is used on hole j can be expressed as Eq. (1):

$$Y_{mj} = \frac{t_{mj}}{T_{mj}} \times Z_m + Yt_{mj}, \quad (1)$$

where,

m , tool type index in ascending order according to the tool diameters, $m=1, \dots, M$

j, k, l , hole index, $j=1, \dots, J$ $k=1, \dots, J$ $l=1, \dots, J$

m_j , index for the last tool to be used on hole j

Y_{mj} , combined tool and machining costs when tool type m is used on hole j .

t_{mj} , machining time required by tool m for hole j

T_{mj} life of tool type m associated with cutting operation on hole j

Z_m , cost of tool type m

Y , machining cost per unit time

Machining time, t_{mj} , is determined by Eq. (2):

$$t_{mj} = \frac{\pi d_m L_j}{1000 U_{mj} f_m}, \quad (2)$$

where,

d_m , diameter of tool m

L_j , depth of hole j , including the clearance

U_{mj} , cutting speed of tool m associated with an operation on hole j

f_m , recommended feed rate for tool type m

In drilling operations depth of cut is fixed. In normal practice feed is kept as a constant rate of cutting speed. Hence the optimum cutting speed, U_{mj} , for the constant feed rate can be obtained by solving the following differential Eq. (3):

$$\frac{dY_{mj}}{dU_{mj}} = 0 \quad (3)$$

The cutting speed obtained from Eq. (3) reduces the sum of tool and machining costs for a single operation. Considering all aspects mentioned above the final optimization model can be expressed as given by Eqs. (4-6).

$$\min G(s) = \min \sum_{m \in M_j} \sum_{m' \in M_j} \sum_{m'' \in M_j} \sum_{l=1}^J \sum_{j=1}^J \sum_{k=1}^J x_{mm'm''ljk} \left[a \left(\frac{p_{lj} + p_{jk}}{2} \right) + b q_{mm'j} + \frac{t_{mm'j}}{T_{mm'j}} Z_m + t_{mm'j} Y \right] \quad (4)$$

subject to

$$\sum_{m \in M_j} \sum_{m' \in M_j} \sum_{l=1}^J \sum_{k=1}^J x_{mm'm''ljk} = 1 \forall j \quad (5)$$

$$x_{mm'm''ljk} + x_{mm'm''kjl} \leq 1 \{ l, j, k, m, m', m'' \} l \neq j, k \neq j, m \in M_j, m' \in M_j, m'' \in M_j \quad (6)$$

where,

s is the sequence index, denoting a specific permutation of operations, $G(s)$ represents the overall cost associated with operations in sequence s , a denotes the cost per unit non-productive travelling distance, b is associated with the cost per unit tool switch time, M_j is a set of tools that can be used to drill hole j to its final size, p_{jk} , is non-productive travelling distance between hole j and hole k , $q_{mm'j}$, represents tool switch time between current tool type, m'' , and tool m required by hole j and finally $x_{mm'm''ljk}$, is a 0-1

integer variable, i.e. $x_{mm'lj} = 1$ if tool m replaces tool m' to drill hole j which is located in the path between holes l and k and has been drilled by tool m' ; 0, otherwise. The 0-1 decision variable, $x_{mm'lj}$ simultaneously determines the sequence of holes to be processed as well as the indices m , m' , and m'' are used to achieve the correct sequence of tools for machining of each hole. Constraint set (5) ensures that each hole is drilled to its final size. Constraint set (6) states that backward movement of spindle is not allowed unless a tool switch is needed. To solve this model large amount of computational time is required as relatively large number of 0-1 decision variables are involved. To overcome this problem, efficient solution procedure using PSO algorithm is proposed.

3. Particle Swarm Optimization (PSO) algorithm

Particle swarm optimization is an evolutionary computation technique developed by Kennedy and Eberhart (1995). The particle swarm thought was originated as a simulation of a simplified social system. This technique starts with initialization of population of random solutions called “particles”. This algorithm consists of two “best” values. First one is the “ p_{best} ” best fitness values of individual particles achieved so far. Second is the “ g_{best} ” which is the one with the best values among all the particles. Velocity and position of individual particles are obtained and updated using Eqs. (7-8) (Kennedy & Eberhart, 1995). Each particle updates its velocity and position through the problem space by comparing its current position and velocity with the optimal solution. In PSO, velocity of particles is changed at every generation towards the “ p_{best} ” and “ g_{best} ”.

$$V_{i+1} = w \times V_i + C_1 \times r_1 \times (p_{besti} - X_i) + C_2 \times r_2 \times (g_{besti} - X_i) \quad (7)$$

$$X_{i+1} = X_i + V_{i+1} \quad (8)$$

where,

V_{i+1} = New velocity of each particle,

w = Inertia weight,

V_i = Previous velocity of particle,

r_1 & r_2 = random numbers between 0 to 1,

C_1 & C_2 = acceleration constants or Cognitive and social constants,

X_i = Previous position of particle.

The acceleration constants ‘ C_1 ’ and ‘ C_2 ’ in Eq. (7) represent the weighting of the stochastic acceleration terms that pull each particle towards ‘ p_{best} ’ and ‘ g_{best} ’ positions. Thus, tuning of these constants varies the amount of tension in the system. Low values of the constants allow particles to pass through far from target regions before being tugged back, while high values result in rapid movement toward, or pass through target regions (Elbeltagi et al. 2005; Dong et al. 2005). The inertia weight ‘ w ’ plays a crucial role in the PSO convergence behavior since it is used to manage the exploration abilities of the swarm. The effect of w , C_1 and C_2 on convergence for standard numerical benchmark functions was provided by Bergh and Engelbrecht (2006). The optimum selection of operating parameters of the algorithm like acceleration constants ‘ C_1 ’ and ‘ C_2 ’ as well as inertia weight ‘ w ’ is essential for the convergence of the algorithm. To ensure the convergence of the PSO algorithm, the condition specified by the f Eq. (9) must be satisfied (Bergh & Engelbrecht, 2006):

$$\max (|\lambda_1|, |\lambda_2|) < 1 \quad (9)$$

where

$$\lambda_1 = \frac{1 + w - \phi_1 - \phi_2 + \gamma}{2}, \tag{10}$$

$$\lambda_2 = \frac{1 + w - \phi_1 - \phi_2 - \gamma}{2}, \tag{11}$$

where $\gamma = \sqrt{(1 + w - \phi_1 - \phi_2)^2 - 4w}$, $\phi_1 = C_1 \times r_1$ and $\phi_2 = C_2 \times r_2$.

As the feasible range for w is 0-1 and for C_1 and C_2 is 0-2, the selected values of w , C_1 and C_2 should be arranged such that the Eq. (9) is satisfied for all possible values of random numbers r_1 and r_2 in the range 0-1. The controlling parameters of PSO algorithm are selected based on the above mentioned criteria for the application example discussed in the next section.

4. Application Example

The particle swarm optimization algorithm is implemented to determine the optimal sequence of operations and corresponding cutting speeds of the upper holder of plastic injection mould as shown in Fig. 2 (Kolahan & Liang, 2000). The input data required for determining the optimal sequence of operations and corresponding cutting speeds of this mould using PSO are considered from (Kolahan & Liang, 2000). This mould consists of total 32 holes namely GP1, GP2, GP3, GP4, GE1-GE4, PR1-PR4, C1-C4, C1"-C4", P1-P4, EB1-EB6, ES1-ES2. Fig. 2 also shows data related to the distances between the holes, type of operations required, and the depth of each hole.

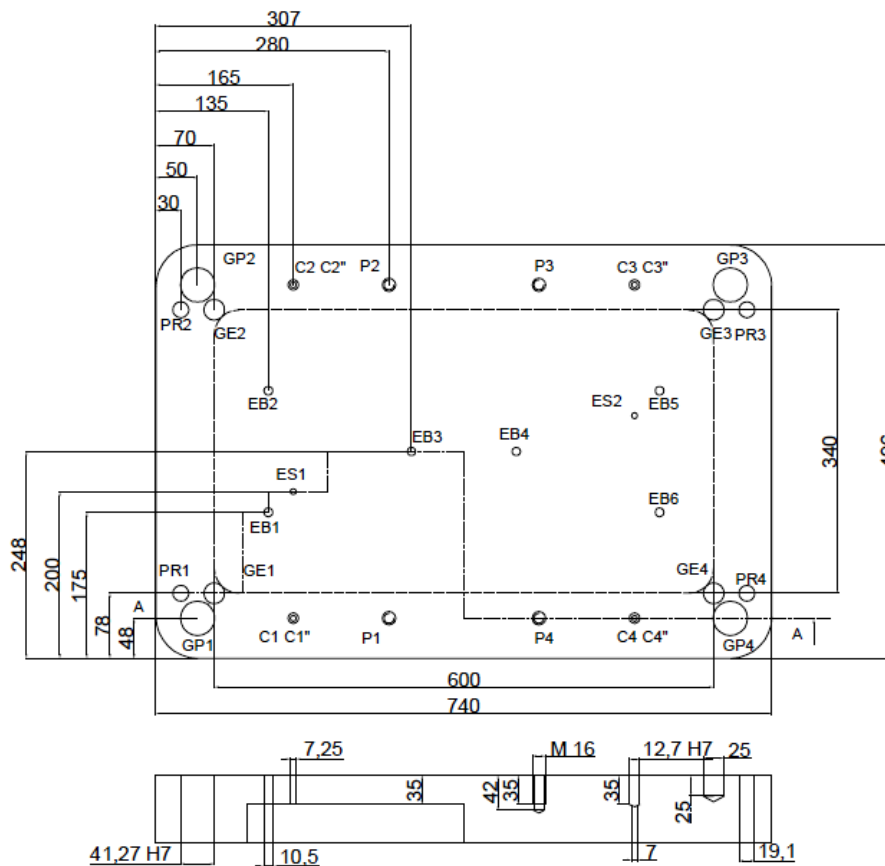


Fig. 2. Upper holder of the plastic injection mould

Three types of operations: drilling, reaming, or tapping are necessary to machine the holes on this part. Total numbers of tools required for hole-making are 12. Data of each tool and its corresponding feed rate, diameter and cost of machining are given in Table 1. The tool life expressions for drilling, reaming, or tapping operations are given in Eqs. (12- 14) (Zhao, 1992).

Table 1

Data of tool diameter, cost, and specified feed rate considered for an application example.

Tool type <i>m</i>	Drill							Reamer			Tap	
	1	2	3	4	5	6	7	8	9	10	11	12
<i>f_m</i> (mm/rev)	0.12	0.1	0.12	0.15	0.2	0.2	0.18	0.15	0.5	0.8	0.8	1.5
<i>d_m</i> (mm)	7	7.25	10.5	12.5	13	19	25	41	12.7	19.1	41.2	16
<i>Z_m</i> (\$)	10	12	15	15	14	20	26	50	55	70	85	45

$$T_{mj} = \left(\frac{8d_m^{0.4}}{U_{mj} \times f_m^{0.7}}\right)^5, \quad \text{for drilling a new hole} \tag{12}$$

$$T_{mj} = \left(\frac{18.4d_m^{0.4}}{U_{mj} \times e_{mj}^{0.2} \times f_m^{0.7}}\right)^5, \quad \text{for enlarging a hole by drilling} \tag{13}$$

$$T_{mj} = \left(\frac{12.1d_m^{0.3}}{U_{mj} \times e_{mj}^{0.2} \times f_m^{0.65}}\right)^{2.5}, \quad \text{for enlarging a hole by reaming or tapping} \tag{14}$$

where, *e_{mj}*, is depth of cut when tool type *m* performing a cutting operation on hole *j*. Optimum cutting speeds expressed in Eqs.(15-17) can be achieved by solving differential Eq.(3) with Eqs. (12-14) (Kolahan & Liang, 2000):

$$U_{mj} = 6 \times 5 \sqrt{\frac{Yd_m^2}{Z_m f_m^{3.5}}}, \quad \text{for drilling a new hole} \tag{15}$$

$$U_{mj} = 13.9 \times 5 \sqrt{\frac{Yd_m^2}{Z_m e_{mj} f_m^{2.5}}}, \quad \text{for enlarging a hole by drilling} \tag{16}$$

$$U_{mj} = 10.3 \times 2.5 \sqrt{\frac{Yd_m^{0.75}}{Z_m e_{mj}^{0.5} f_m^{1.65}}}, \quad \text{for enlarging a hole by reaming or tapping} \tag{17}$$

Tooling and machining costs of individual operations are calculated using optimum cutting speeds obtained using Eqs. (15-17). For this application example, tool switch times required for machining of hole-making operations are given in Table 2.

Table 2

Tool switch times (min)

Next in line Tool	Previous Tool											
	1	2	3	4	5	6	7	8	9	10	11	12
1	0	0.6	0.2	0.4	0.4	0.9	0.6	1	0.8	1.4	0.4	1
2	0.6	0	0.8	1.2	0.4	0.8	1.2	0.5	0.6	0.6	1.2	1.4
3	0.2	0.8	0	0.6	1.4	1.2	1.4	1	0.4	1	1.9	1.2
4	0.4	1.2	0.6	0	0.4	0.7	1	1.6	0.8	1.4	0.4	0.6
5	0.4	0.4	1.3	0.5	0	0.8	0.6	1	0.8	1.8	1.6	1.5
6	0.5	0.5	1.2	0.2	0.8	0	0.2	1.5	1.2	1.3	0.4	1.9
7	0.6	1.2	1.4	1	0.6	0.4	0	1	0.8	0.8	1.2	0.6
8	1	0.6	0.6	0.4	0.4	0.7	0.6	0	0.8	1.5	0.4	0.5
9	0.2	0.9	1.1	0.6	1.5	1.2	1.4	1	0	1.2	2	1.2
10	0.4	1.2	0.7	1.5	0.5	0.3	1	0.4	0.7	0	0.4	0.6
11	0.4	1.2	2	0.4	1.2	0.8	0.6	1.2	0.8	1	0	0.2
12	1	0.7	0.3	0.4	0.4	0.8	0.6	1.4	1	1.5	0.5	0

Table 3

Tool-hole combinations considered in the application example.

J	GP1-GP4	GE1-GE4	PR1-PR4	C1-C4	C1"-C4"	P1-P4	EB1-EB6	ES1-ES2
m_j	6-8-11	6-7	6-10	4-9	1	5-12	3	2

Table 3 corresponds to different combinations of tools necessary for machining of an individual hole to its final size as shown in Fig. 2. For example, holes GP1, GP2, GP3, GP4 require tool number 6 as initial tool, tool number 7 as intermediate tool and tool number 11 and a reamer to achieve the final size hole. Similarly, other holes involve a tool or different combinations of tools to achieve the final size hole. As given in Table 3, 56 machining operations are required for the example shown in Fig. 2. Process parameter assumed for this application example are, $Y=\$1/\text{min}$, $a=\$0.0008/\text{mm}$ and $b=\$1/\text{min}$.

5. Results and Discussion

For computational experiments, a Windows 8 PC with Intel core i3 CPU @ 1.90 GHz and Code blocks C compiler were used. In order to compare the results of PSO with those obtained using tabu search method developed by Kolahan and Liang (2000), for the application example considered in section 4, the following two cases are considered:

Case 1: Considering tool switch times given in Table 2,

Case 2: Considering tool switch times 50% of those given in Table 2,

Case 1: Following algorithm specific parameters for particle swarm optimization are obtained through various computational experiments.

$C_1=1.5$,

$C_2=2.0$,

$w=0.5$,

Number of iterations =500,

Number of particles=100.

For the above parameter setting, the results of optimization for case 1 using PSO are given in Table 4.

Table 4

Results of optimal sequence of operation and associated cutting speeds for Case 1 using PSO

m_j	6-GE2	6-GP2	6-PR2	6-PR1	6-GP1	6-GE1	6-GE3	6-GP3	6-PR3
U_{mj}	33.016	33.016	33.016	33.016	33.016	33.016	33.016	33.016	33.016
m_j	6-PR4	6-GE4	6-GP4	8-GP4	8-GP3	8-GP2	8-GP1	11-GP1	11-GP2
U_{mj}	33.016	33.016	33.016	44.876	44.876	44.876	44.876	9.761	9.761
m_j	11-GP3	10-PR3	10-PR4	10-PR1	10-PR2	7-GE3	7-GE2	7-GE1	1-C4"
U_{mj}	9.761	9.622	9.622	9.622	9.622	49.675	49.675	49.675	36.372
m_j	4-C4	9-C4	5-P4	5-P1	4-C1	4-C2	1-C2"	1-C1"	9-C1
U_{mj}	36.177	11.13	30.464	30.464	36.177	36.177	36.372	36.372	11.13
m_j	9-C2	5-P2	5-P3	12-P4	12-P3	12-P2	12-P1	7-GE4	11-GP4
U_{mj}	11.13	30.464	30.464	3.642	3.642	3.642	3.642	49.675	9.761
m_j	4-C3	1-C3"	9-C3	3-EB5	3-EB1	3-EB3	3-EB2	3-EB4	3-EB6
U_{mj}	36.177	36.372	11.13	39.444	39.444	39.444	39.444	39.444	39.444
m_j	2-ES2	2-ES1							
U_{mj}	40.406	40.406							

Table 4 corresponds to the optimal sequence of operations and associated cutting speeds of Case 1 that are obtained using PSO. This sequence results into overall processing cost of \$66.78 including \$45.2 machining and tool costs, \$10.48 non-productive travelling cost, and \$11.1 tool switch cost.

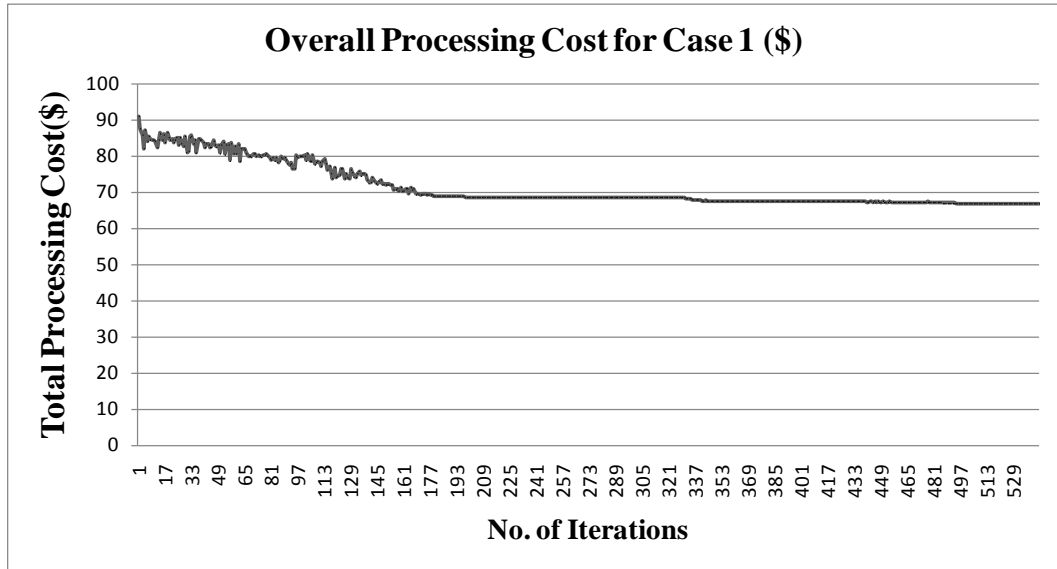


Fig. 3.a. Convergence of overall processing costs (\$) of Case 1 using PSO

Case 2: The following algorithm specific parameters for particle swarm optimization are obtained through various computational experiments.

- $C_1=1.65,$
- $C_2=1.75,$
- $w=0.65,$
- Number of iterations =600,
- Number of particles=100.

Table 5

Results of optimal sequence of operation and associated cutting speeds of Case 2 using PSO

m_j	3-EB2	3-EB3	3-EB4	3-EB6	3-EB5	2-ES2	2-ES1	6-PR1	6-GP1
U_{mj}	39.444	39.444	39.444	39.444	39.444	40.406	40.406	33.016	33.016
m_j	6-GE1	7-GE1	8-GP1	10-PR1	11-GP1	4-C1	1-C1"	9-C1	5-P1
U_{mj}	33.016	49.675	44.876	9.622	9.761	36.177	36.372	11.13	30.464
m_j	5-P4	4-C4	6-GP4	8-GP4	1-C4"	9-C4	12-P4	11-GP4	6-GE4
U_{mj}	30.464	36.177	33.016	44.876	36.372	11.13	3.642	9.761	33.016
m_j	7-GE4	6-GE3	6-PR3	6-PR4	10-PR4	10-PR3	7-GE3	6-GP3	8-GP3
U_{mj}	49.675	33.016	33.016	33.016	9.622	9.622	49.675	33.016	44.876
m_j	11-GP3	4-C3	1-C3"	9-C3	5-P3	5-P2	12-P3	12-P2	4-C2
U_{mj}	9.761	36.177	36.372	11.13	30.464	30.464	3.642	3.642	36.177
m_j	1-C2"	9-C2	6-GP2	6-GE2	6-PR2	7-GE2	8-GP2	11-GP2	10-PR2
U_{mj}	36.372	11.13	33.016	33.016	33.016	49.675	44.876	9.761	9.622
m_j	3-EB1	12-P1							
U_{mj}	39.444	3.642							

Table 5 corresponds to the optimal sequence of operations and associated cutting speeds of Case 2 that are obtained using PSO. This sequence results into an overall processing cost of \$60.45 from which \$45.2 is the tool cost and machining cost, \$10.94 tool switch cost, and \$4.31 tool travel cost.

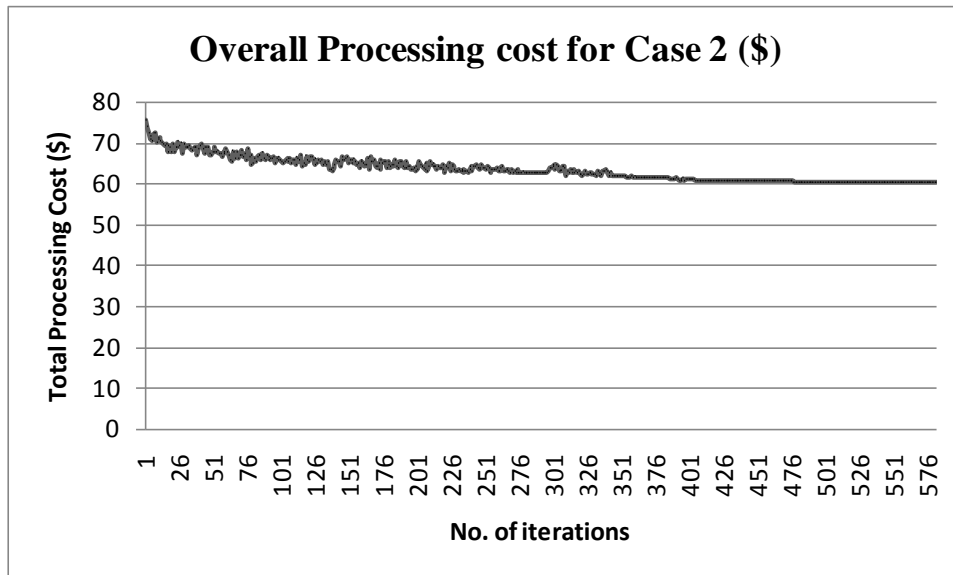


Fig. 3.b. Convergence of overall processing costs (\$) of Case 2 using PSO

Table 6 shows comparison of results of application example obtained by using PSO algorithm and tabu search (Kolahan & Liang 2000).

Table 6

Comparison of results of optimization obtained by using PSO with those obtained by using tabu search Kolahan and Liang (2000) for Case 1 and Case 2

Case 1	Tooling and Machining Cost C_{mj} (\$)	Tool Travel Cost (\$)	Tool Switch Cost (\$)	Overall Processing Cost (\$)
Tabu Search (Kolahan and Liang2000)	45.2	11	8.6*	64.8
Tabu Search (Kolahan and Liang2000)	45.2	11	11.2**	67.4
PSO	45.2	10.48	11.1	66.78
Case 2				
Tabu Search (Kolahan and Liang 2000)	45.2	4.9	10.1*	60.2
Tabu Search (Kolahan and Liang 2000)	45.2	4.9	13.15**	63.25
PSO	45.2	4.31	10.94	60.45

*Value wrongly calculated by Kolahan and Liang (2000).

**Corrected values obtained by substituting the optimum result obtained by Kolahan and Liang (2000) in Eq. (4)

The example of this application was originally solved by Kolahan and Liang (2000) using tabu-search approach in order to reduce the overall processing cost of hole-making operations. Sequence obtained using tabu-search for both cases is checked manually as given Eq. (4), it is observed that the actual tool switch cost for both cases is different than the results given by Kolahan and Liang (2000). Corrected results for both cases are given in Table 6. PSO results are compared with these corrected results given in Table 6.

6. Conclusion

Optimization of hole-making operations involves large number of hole-making operations sequences due to the location of hole and tool sequence constraint. To achieve this, proper determination operations sequence and associated cutting speeds which reduces the overall processing cost of hole-making operations are essential. In this paper, a methodology has been proposed to reduce the overall processing cost of hole-making operations of an application example using PSO algorithm. The obtained results

have been compared with those obtained using tabu-search approach reported by Kolahan and Liang (2000). It is observed that the results of optimization obtained by PSO algorithm were slightly better than tabu-search approach (Kolahan & Liang, 2000) since for both cases showing an improvement about 1.0% for Case 1 and 4.6 % for Case 2. However for the both cases, the sequence of operation to be performed shows significant changes with respect to results obtained using tabu-search approach (Kolahan & Liang 2000). This clearly shows that PSO algorithm has potential to solve this problem. Also it is observed that PSO algorithm requires only 600 generations to converge to optimal solution. The improvement obtained by using PSO algorithm is thus significant and clearly indicates the potential of this method to solve real life problems related to hole-making for various industrial applications.

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