

Harmony search algorithm with adaptive parameter setting for solving large bin packing problems

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ABSTRACT

Bin packing problem is a constrained optimization problem with a huge search space due to large combinations. Bin packing problem has a wide range of applications in multiple fields. This paper presents harmony search algorithm with different initialization and adaptive PAR strategies for solving bin packing problem. The proposed Harmony search (HS) variations tests two partial feasible initialization strategies for bin packing problem. The paper presents adaptive PAR strategies for better exploration and exploitation of HS algorithm. The PAR values are tuned in every iteration. Improved initialization strategy, population initialization after premature convergence and adaptive PAR leads to the better exploration of harmony search algorithm for bin packing problem. The performance of variations are tested over 120 benchmark instances with 100 and 200 objects with varying complexities. The results show that improved HS performs better than basic HS with respect to best, mean, convergence rate. The performance of algorithms is tested with varying harmony memory size and harmony memory considering rate. Results show that variation in these two parameter values has less effect on performance of improved versions.

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1. Introduction

Combinatorial optimization problems are difficult to solve due to huge search space and constraints. Bin packing problem (BPP) is one of the important combinatorial optimization problems. Bin packing problem has applications in multiple domains such as transportation (Perboli et al., 2014), logistics (Aggoun et al., 2016), resource placement (Song et al., 2013; Adamuthe & Patil, 2018). The BPP problem formulation is well known and presented in many papers (Fleszar & Hindi, 2002; Schoenfeld, 2002; Alvim et al., 2004; Abdel-Basset et al., 2018). In one-dimensional bin packing problem, items with varying weights are to be packed using homogeneous/heterogeneous bins. Consider a set of 'k' bins $Bin_1, Bin_2 \dots Bin_k$ with same or different capacity and 'n' objects with different weights $W_1, W_2 \dots W_n$. The objective is to pack all the items into the bins in such a way that a minimum number of bins are required. A solution is said to be optimal if it uses a minimal number of bins.

$$\min \sum_{i=1}^k z_i \quad (1)$$

subject to

$$\sum_{j=1}^n W_j y_{ij} \leq B_i \quad \text{for all } i \quad (2)$$

$y_{ij} = 1$ if object j is put in bin i , zero, otherwise

$z_i = 1$ if bin i is used, 0, otherwise

$i \in k$ and $j \in n$

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In literature, bin packing problem is investigated with different deterministic and heuristic algorithms. First-Fit (FF) and Best-Fit (BF) algorithms have an absolute worst-case ratio of 7/4 for the BPP (Simchi-Levi, 1994). Worst-case performance ratio of first fit algorithm (Dósa & Sgall, 2013) and best fit algorithm is 17/10 (Dósa & Sgall, 2014). A large investigation is done to solve optimization problems using approximation algorithms. Approximation algorithms are investigated when polynomial-time of an algorithm is expensive. A linear approximation algorithm was presented in (Berghammer & Reuter, 2003) for BPP with absolute approximation factor 3/2. Martel (1985) presented a linear off-line algorithm for bin packing with performance ratio of 4/3. In literature, different heuristic algorithms are experimented to solve bin packing problem such as genetic algorithms (Falkenauer, 1996), cuckoo search (Layeb & Boussalia, 2012; Zendaoui & Layeb, 2016), firefly algorithm (Layeb & Benayad, 2014), ant colony optimization (Levine & Ducatelle, 2004), Lévy-based whale optimization algorithm (Abdel-Basset et al., 2018).

In 2001, Geem et al. (2001) presented a new population based heuristic algorithm called as Harmony Search (HS) algorithm. The improvisation technique used by HS algorithm is similar to that of music players. HS is analogous to GA and PSO. In recent years harmony search algorithm is investigated to solve different optimization problems (Lee & Geem, 2005; Mahdavi et al., 2007; Kattan & Abdullah, 2013; Valian et al., 2014). From 2001, various modification have been proposed by researchers to improve efficiency of basic HS algorithm. These modified HS algorithms are better than basic HS in terms of convergence, optimization capability and solution accuracy for applied problem instances. An opposition-based initialization to accelerate HS algorithm was proposed by Banerjee et al. (2014) and results conclude robustness of this approach. Taherinejad (2009) presented modified harmony search (MHS) with pitch-adjustment step in such a way that harmony is affected by the best harmony in the HM. Different investigation are,

- Pitch adjustment rule (Omran & Mahdavi, 2008)
- Balanced intensification and diversification with suitable pitch adjustment rules (Yadav et al., 2012)
- Adaptive parameter setting (Mahdavi et al., 2007; Wang & Huang, 2010; Khalili et al., 2014; Park et al., 2017)

This paper presents harmony search algorithm for solving bin packing problem. The basic harmony search algorithm shows good convergence towards the optimal solution. However, stuck in performing local search in the high performance region in the search space. Paper presents partially feasible initialization, re-initialization strategies and adaptive PAR for bin packing problem. Paper presents two strategies for changing PAR values

- Increasing PAR presented by (Mahdavi et al., 2007)
- Decreasing PAR

Performance of basic HS and two variations are presented using 120 benchmark instances of bin packing problem with 100 and 200 objects. Experiments are conducted to check the effect of Harmony Memory Size (HMS) and Harmony Memory Considering Rate (HMCR).

In next section proposed harmony search algorithm for solving bin packing problem is presented. Section 3 is about dataset, results and performance evaluations. At last, conclusions is stated in section 4.

2. Harmony search for Bin packing problem

Harmony search is inspired from musicians' behaviour. Musicians' try to improvise pitches of their instruments. The fundamental steps of harmony search algorithm are (Geem et al., 2001):

- Step 1. Harmony search algorithm parameters initialization
- Step 2. Harmony memory initialization
- Step 3. Improvise harmony
- Step 4. Update the harmony memory
- Step 5. Repeat steps 3 and 4 until the termination condition is reached

This section presents optimizing bin packing problem with harmony search algorithm. Application of basic steps of harmony search algorithm for bin packing are explained in detail. Bin packing problem is formulated as minimization problem in which we have to minimize the number of bins used. Penalty function is applied if constraint described by equation 3 is violated. The total penalty for a solution is calculated as follows:

$$P = \left| \sum_{i=1}^k p_i \right| \quad (3)$$

Penalty value of the solution is considered as objective value.

Step 1: This is basic step of HS. Need to initialize the HS parameters namely harmony memory size (HMS), harmony memory consideration rate (HMCR), pitch adjusting rate (PAR). Paper presents basic and two version of adaptive HS

algorithm. For basic HS algorithm, parameters are initialized randomly. For adaptive PAR HS algorithm the initialization values for PAR are presented below.

Step 2: Representation and initialization of harmony memory (HM) considering the problem specific knowledge.

Real coded solution representation is used to solve the BPP problem. Size of the array depends on the number of objects to be packed in bin. The values in the array indicate the bin number to which the object is assigned. Fig. 1 shows the sample solution with 10 objects and 6 bins. For example, Obj₅ is assigned to Bin₁. Objects Obj₂ and Obj₅ are packed in bin Bin₁. This representation guarantees the placement of all objects.

Obj ₁	Obj ₂	Obj ₃	Obj ₄	Obj ₅	Obj ₆	Obj ₇	Obj ₈	Obj ₉	Obj ₁₀
Bin ₃	Bin ₁	Bin ₅	Bin ₆	Bin ₁	Bin ₂	Bin ₅	Bin ₉	Bin ₃	Bin ₂

Fig. 1. Solution Representation

In HS a matrix known as harmony memory (HM) which stores best solution vectors is considered as core data structure. HM can be represented as a matrix where each row corresponds to solution vector. Sample HM is shown in Fig 2.

Bin ₃	Bin ₁	Bin ₉	Bin ₆	Bin ₁	Bin ₂	Bin ₅	Bin ₉	Bin ₃	Bin ₈
Bin ₉	Bin ₃	Bin ₂	Bin ₁	Bin ₅	Bin ₉	Bin ₆	Bin ₁	Bin ₂	Bin ₂
Bin ₃	Bin ₁	Bin ₅	Bin ₆	Bin ₁	Bin ₂	Bin ₅	Bin ₉	Bin ₃	Bin ₈
Bin ₈	Bin ₅	Bin ₉	Bin ₆	Bin ₁	Bin ₂	Bin ₅	Bin ₉	Bin ₃	Bin ₃
Bin ₅	Bin ₂	Bin ₉	Bin ₆	Bin ₃	Bin ₃	Bin ₉	Bin ₁	Bin ₇	Bin ₁

Fig. 2. Sample harmony memory

This paper presents two problem specific harmony search initialization methods for solving bin packing problem.

HM initialization procedure 1:

In this initialization strategy, we randomly select a bin for all object sequentially. The randomly selected bin is assigned to object if it has capacity to occupy the object i.e. (current_bin_capacity – obj_size) > 0. If the assignment violates the constraint then next fit bin is identified.

HM initialization procedure 2:

In this initialization strategy we first iterate over all the bins ($\forall i. 0 \leq ii \leq \text{num_bins}$) until we find a bin such that $\text{bin_size}[ii] - \text{obj_size} \geq 0$ and then we assign that bin to current object. If after iterating all bins, we fail to find a bin which can accommodate a object then, we randomly assign a bin to that object.

```

reference_bin[ii][jj] ← bin_cap,  $\forall i. ii \leq \text{HMS} \wedge \forall j. jj \leq \text{number\_of\_bins}$ 
while ii ≤ HMS do
  while jj ≤ number_of_obj do
    assigned_bin ← rand[0, max_bins]
    if obj_weight[jj] < reference_bin[ii][assigned_bin] then
      HM[ii][jj] ← assigned_bin
      reference_bin[ii][jj] ← reference_bin[ii][jj] – obj_weight[ii]
    else
      while kk ≤ number_of_bins do
        if obj_weight[jj] < reference_bin[ii][kk] then
          HM[ii][jj] ← kk
          reference_bin[ii][kk] ← reference_bin[ii][kk] – obj_weight[ii]
        fi
      od
    fi
  od
od
od

```

Fig. 3. Pseudocode for harmony memory initialization procedure 1

```

reference_bin[ii][jj] ← bin_cap, ∀i. ii ≤ HMS ∧ ∀j. jj ≤ number_of_bins
while ii ≤ HMS do
  while jj ≤ number_of_obj do
    obj_assigned ← False
    while kk ≤ number_of_bins do
      if obj_weight[jj] < reference_bin[ii][assigned_bin] then
        HM[ii][jj] ← assigned_bin
        reference_bin[ii][jj] ← reference_bin[ii][jj] - obj_weight[ii]
        obj_assigned ← True
        break
      fi
    od
    if obj_assigned not True:
      bin ← randomly assign a bin
      HM[ii][jj] ← bin
      reference_bin[ii][bin] ← reference_bin[ii][bin] - obj_weight[ii]
    fi
  od
od

```

Fig. 4. Pseudocode for harmony memory initialization procedure 2

Step 3: Improve a new solution from HM.

Many heuristic algorithms use parameters to control the exploration and exploitation capabilities of algorithm. In this step new solution is generated using two parameters namely harmony memory consideration rate (HMCR) and pitch adjusting rate (PAR). The value of the parameter HMCR and PAR range from 0 to 1. The HMCR is probability of selecting a value from previous HM. Every selected value by the memory consideration is examined is modified considering pitch adjusting rate. HMCR and PAR are important for adjusting the convergence of algorithm and fine tune the results to get the optimized solution. Basic harmony search uses static values in all iterations until stopping criteria is satisfied. This paper presents two ways for adaptive PAR values based on the convergence rate of algorithm. The value of PAR is changed after every iteration depending on HM diversity.

Adaptive PAR strategy 1:

The value of PAR is tuned as per following equation (Mahdavi et al., 2007):

$$PAR = PAR + (PAR_{min} + \frac{(PAR_{max} - PAR_{min})}{NI} \times gn) \quad (4)$$

gn is the current generation (current iteration) and NI represents maximum iterations. PAR_{max} and PAR_{min} are initialized to 1.0 and 0.45 respectively.

Adaptive PAR strategy 2:

$$PAR = PAR + (PAR_{min} - \frac{(PAR_{max} - PAR_{min})}{NI} \times gn) \quad (5)$$

PAR_{max} and PAR_{min} are initialized to 1.0 and 0.85 respectively.

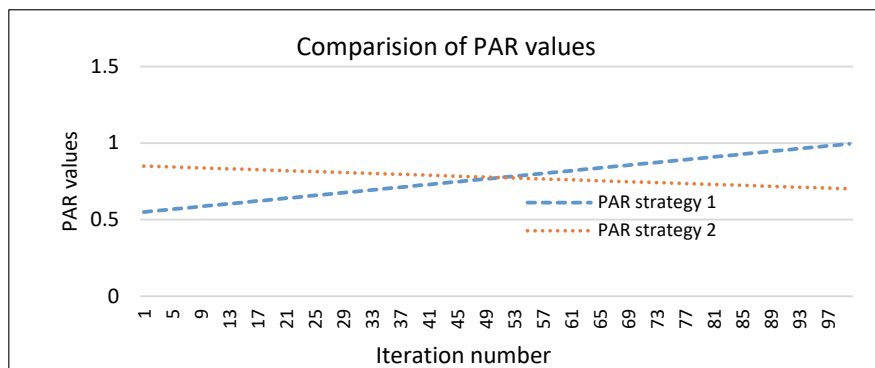


Fig. 5. Change in PAR with iteration numbers

Fig. 6 presents pseudocode for basic harmony search algorithm for solving bin packing problem. Fig. 7 presents pseudocode for adaptive harmony search algorithm with dynamic PAR. It covers two versions of adaptive HS algorithm with different PAR initiation and changing strategy.

Basic Algorithm (V1): Harmony Search with static PAR

Input: Number of objects with weight, number of bins with capacity
Variables: number_of_bins, bin_cap, HS, number_of_obj, assigned_bin, obj_weight
Output: minimum bins required

Initialize HMS, HMCR, PAR and max iterations
 Initialize harmony memory with procedure 1

```

begin
  while  $ii \leq \text{max\_iterations}$  do
    while  $jj \leq \text{HMS}$  do
      while  $kk \leq \text{number\_of\_objects}$  do
        if  $\text{rand}[0,1] < \text{HMCR}$  then
          memory_consideration( $kk$ )
          if  $\text{rand}[0,1] < \text{PAR}$  then pitch_adjustment( $jj, kk$ )fi
        else
          random_solution( $jj, kk$ )
        fi
      od
      memory_updation( $jj$ )
    od
  od
end

```

Fig. 6. Pseudocode Harmony Search with static PAR (V1)

Adaptive Harmony Search with Dynamic PAR

Input, Variables, Output: ← same as Algorithm 1

Initialize HMS, HMCR and max iterations
 Initialize PAR with strategy 1 (for version 2 – V2)
 Initialize PAR with strategy 2 (for version 3 – V3)
 Initialize harmony memory with procedure 2

```

begin
  while  $ii \leq \text{max\_iterations}$  do
    while  $jj \leq \text{HMS}$  do
      while  $kk \leq \text{number\_of\_objects}$  do
        if  $\text{rand}[0,1] < \text{HMCR}$  then
          memory_consideration( $kk$ )
          if  $\text{rand}[0,1] < \text{PAR}$  then
            if  $\text{reference\_bin}[jj][kk] > \text{bin\_cap}$  then
              pitch_adjustment( $jj, kk$ )
            fi
          fi
        else
          random_solution( $jj, kk$ )
        fi
      od
      memory_updation( $jj$ )
    od
  od

  Adaptive PAR strategy 1 (for variation 2- V2)
  Adaptive PAR strategy 2 (for variation 3- V3)

  if  $\text{current\_iteration} > \text{threshold}$  then
    re_initialize_HM()
    threshold = threshold + 100
  fi
od
end

```

Fig. 7. Pseudocode Adaptive Harmony Search with dynamic PAR (V2 and V3)

3. Results and Performance Evaluation of HS

This section presents dataset details, obtained results and discussion. The proposed harmony search algorithm variations for solving bin packing problem are implemented using ‘C’ programming language. The programs are executed on machine with Intel core i5-3210M CPU 2.5 GHz and 4 GB RAM. Benchmark instances of bin packing problem are taken to test the performance of proposed variations. Benchmark datasets for bin packing problem taken from <https://www2.wiwi.uni-jena.de/Entscheidung/binpp/index.htm>. These data sets are prepared as per Martello and Toth (Martello & Toth, 1990)

and best known solutions are available. The naming conventions of the data sets are as follows.

Parameters used,

- $o \rightarrow$ number of objects
- $c_bin \rightarrow$ capacity of bin
- $W_j \rightarrow$ size of object j ($j=1, \dots, o$)
- $min \rightarrow$ minimum number of required bins

The name of the instances are given by $NpCqWr_i$ where,

- $p=1$ (for $o=50$), $p=2$ (for $o=100$), $p=3$ (for $o=200$), $p=4$ (for $o=500$)
- $q=1$ (for $c_bin=100$), $q=2$ (for $c_bin=120$), $q=3$ (for $c_bin=150$)
- $r=1$ (for W_j from $[1,100]$), $r=2$ ($[20,100]$), $r=4$ ($[30,100]$, for ($j=1, \dots, o$))
- $i=A \dots T$ for the 20 instances of each class

Total 120 instances with $p=2$ and 3, $q=1$ and $r=1, 2$, and 4 are used for experimentation. The input file for the following instances consists of number of objects, bin capacity and size of each object separated by new line. The structure of input data sets is as given below.

Sample input file	
3	← number of objects
100	← capacity of bin
30	} size of items
40	
30	

These problem instances are experimented by different authors. The optimal solutions are presented by (Fleszar & Hindi, 2002; Schoenfeld, 2002; Alvim et al., 2004; Schwerin & Wäscher, 1997). Paper presents the performance evaluation of proposed variations of harmony search algorithms using obtained best, mean, standard deviation, convergence graph and success rate. The success rate is calculated using equation 6. For every dataset harmony search algorithm was executed 10 times.

$$\text{Success rate} = \text{Number of times best obtained} / \text{number of times program executed} \quad (6)$$

Table 1 and 2 shows the success rate of experimented three harmony search algorithm for 120 instances of bin packing problem. The variation V1 fails to give the best known results for all the instances with 100 and 200 objects. Results indicate that proposed initialization strategy and HS algorithm are not tuned to the problem for obtaining best results. Results show that the performance of HS version 2 and version 3 is better than the basic HS algorithm. Both the variations provide best-known solution for given instances. Initialization strategy 2, adaptive PAR and re-initialization works well for bin packing problem. The result shows that adaptive PAR harmony search algorithm versions work better than static PAR HS. Harmony search algorithm with increasing PAR (version 2) is found better than decreasing PAR (version 3). The variation 2 gives 100 percent success rate for more than 66 percent instances with 100 objects. Variation V2 gives good success rate to majority of instances with 100 objects. Approximately 1 percent of instances shows less success rate than 50 percent. The 100 percent success rate is drastically reduced to approximately 25 percent for instances with 200 objects. Approximately 10 percent of instances shows less success rate than 50 percent. There is a large variation on the success rate.

The variation 3 gives 100 percent success rate for more than 66 percent instances with 100 objects. Version 3 fails to give an optimal solution to approximately 25% instances with 100 objects. The performance of version 3 reduces drastically for instances with 200 objects. Only 50% instances are solved by version 3.

Table 1
Success rate of proposed HS variations for 60 instances with 100 objects

Dataset with 100 objects	HS variation			Dataset with 100 objects	HS variation		
	V3	V2	V1		V3	V2	V1
N2C1W1_A	100	100	0	N2C1W2_K	100	100	0
N2C1W1_B	100	100	0	N2C1W2_L	100	50	0
N2C1W1_C	100	80	0	N2C1W2_M	100	100	0
N2C1W1_D	100	100	0	N2C1W2_N	100	90	0
N2C1W1_E	100	100	0	N2C1W2_O	0	100	0
N2C1W1_F	100	100	0	N2C1W2_P	100	100	0
N2C1W1_G	100	90	0	N2C1W2_Q	0	100	0
N2C1W1_H	100	100	0	N2C1W2_R	100	80	0
N2C1W1_I	100	50	0	N2C1W2_S	0	100	0
N2C1W1_J	100	100	0	N2C1W2_T	100	100	0
N2C1W1_K	0	70	0	N2C1W4_A	100	70	0
N2C1W1_L	10	100	0	N2C1W4_B	100	100	0
N2C1W1_M	0	100	0	N2C1W4_C	0	100	0
N2C1W1_N	40	90	0	N2C1W4_D	100	40	0
N2C1W1_O	100	100	0	N2C1W4_E	100	100	0
N2C1W1_P	100	100	0	N2C1W4_F	0	100	0
N2C1W1_Q	0	100	0	N2C1W4_G	0	80	0
N2C1W1_R	0	80	0	N2C1W4_H	60	100	0
N2C1W1_S	100	100	0	N2C1W4_I	100	100	0
N2C1W1_T	100	100	0	N2C1W4_J	100	90	0
N2C1W2_A	100	100	0	N2C1W4_K	0	100	0
N2C1W2_B	0	100	0	N2C1W4_L	100	100	0
N2C1W2_C	100	90	0	N2C1W4_M	100	80	0
N2C1W2_D	100	100	0	N2C1W4_N	100	50	0
N2C1W2_E	0	100	0	N2C1W4_O	100	100	0
N2C1W2_F	100	80	0	N2C1W4_P	80	70	0
N2C1W2_G	100	100	0	N2C1W4_Q	0	100	0
N2C1W2_H	100	100	0	N2C1W4_R	100	100	0
N2C1W2_I	0	100	0	N2C1W4_S	100	100	0
N2C1W2_J	100	60	0	N2C1W4_T	100	100	0

Table 2
Success rate of proposed HS variations for 60 instances with 200 objects

Dataset with 200 objects	HS variation			Dataset with 200 objects	HS variation		
	V3	V2	V1		V3	V2	V1
N3C1W1_A	0	80	0	N3C1W2_K	0	100	0
N3C1W1_B	0	60	0	N3C1W2_L	0	50	0
N3C1W1_C	100	70	0	N3C1W2_M	40	100	0
N3C1W1_D	0	80	0	N3C1W2_N	0	90	0
N3C1W1_E	0	50	0	N3C1W2_O	0	80	0
N3C1W1_F	0	40	0	N3C1W2_P	0	100	0
N3C1W1_G	100	90	0	N3C1W2_Q	20	60	0
N3C1W1_H	0	70	0	N3C1W2_R	0	50	0
N3C1W1_I	0	80	0	N3C1W2_S	100	20	0
N3C1W1_J	100	30	0	N3C1W2_T	0	50	0
N3C1W1_K	0	70	0	N3C1W4_A	0	70	0
N3C1W1_L	0	100	0	N3C1W4_B	0	100	0
N3C1W1_M	100	70	0	N3C1W4_C	0	100	0
N3C1W1_N	0	90	0	N3C1W4_D	0	40	0
N3C1W1_O	0	100	0	N3C1W4_E	0	50	0
N3C1W1_P	0	50	0	N3C1W4_F	0	80	0
N3C1W1_Q	0	100	0	N3C1W4_G	0	80	0
N3C1W1_R	0	80	0	N3C1W4_H	0	60	0
N3C1W1_S	0	40	0	N3C1W4_I	0	60	0
N3C1W1_T	0	80	0	N3C1W4_J	0	80	0
N3C1W2_A	0	40	0	N3C1W4_K	100	70	0
N3C1W2_B	100	100	0	N3C1W4_L	20	60	0
N3C1W2_C	0	90	0	N3C1W4_M	0	80	0
N3C1W2_D	100	100	0	N3C1W4_N	0	50	0
N3C1W2_E	0	100	0	N3C1W4_O	0	100	0
N3C1W2_F	0	80	0	N3C1W4_P	0	70	0
N3C1W2_G	0	100	0	N3C1W4_Q	0	100	0
N3C1W2_H	0	70	0	N3C1W4_R	100	30	0
N3C1W2_I	0	100	0	N3C1W4_S	0	100	0
N3C1W2_J	0	60	0	N3C1W4_T	0	90	0

Table 3 and 4 show the obtained best, mean and standard deviation on the objective value of experimented three harmony search algorithm variations for 120 instances of bin packing problem. The basic harmony search algorithm fails to give best known solution for all the instances. The negative values show the penalty applied for constraint violations. Zero value indicates no violation of constraints. HS variations 2 and 3 shows best performance for more than 95% instances of bin packing problem with 100 objects. The performance of variation 2 and 3 are approximately same.

Table 3

Performance of proposed HS variations for 60 datasets with 100 objects

Dataset	Minimal bins	HS Variation			Dataset	Minimal bins	HS Variation				
		V3	V2	V1			V3	V2	V1		
N2C1W1_A	48	Best	0	0	-285	N2C1W2_K	72	Best	0	0	-313
		SD	0	0	13.56			SD	0	0	14.25
		Mean	0	0	-367.66			Mean	0	0	-398.36
N2C1W1_B	49	Best	0	0	-233	N2C1W2_L	62	Best	0	0	-269
		SD	0	0	18.19			SD	0	0	15.52
		Mean	0	0	-342.45			Mean	0	0	-366
N2C1W1_C	46	Best	0	0	-293	N2C1W2_M	65	Best	0	0	-293
		SD	0	0	13.26			SD	0	0	13.26
		Mean	0	0	-360.45			Mean	0	0	-360.56
N2C1W1_D	50	Best	0	0	-290	N2C1W2_N	64	Best	0	0	-288
		SD	0	0	18.02			SD	0	0	20.92
		Mean	0	0	-370.98			Mean	0	0	-374.38
N2C1W1_E	58	Best	0	0	-318	N2C1W2_O	64	Best	-1	0	-368
		SD	0	0	18.9			SD	2.08	0	14.56
		Mean	0	0	-380.1			Mean	-4	0	-426.36
N2C1W1_F	60	Best	0	0	-328	N2C1W2_P	68	Best	0	0	-384
		SD	0	0	15.42			SD	0	0	21.59
		Mean	0	0	-384.86			Mean	0	0	-498.36
N2C1W1_G	60	Best	0	0	-277	N2C1W2_Q	65	Best	-2	0	-288
		SD	0	0	20.92			SD	1.01	0	20.92
		Mean	0	0	-374.36			Mean	-3	0	-374.36
N2C1W1_H	52	Best	0	0	-369	N2C1W2_R	67	Best	0	0	-532
		SD	0	0	14.86			SD	0	0	21.26
		Mean	0	0	-426.32			Mean	0	0	-609.05
N2C1W1_I	62	Best	0	0	-3842	N2C1W2_S	66	Best	-1	0	-325
		SD	0	0	21.49			SD	0	0	15.36
		Mean	0	0	-467.24			Mean	-1	0	-385.36
N2C1W1_J	59	Best	0	0	-308	N2C1W2_T	66	Best	0	0	-318
		SD	0	0	21.72			SD	0	0	18.9
		Mean	0	0	-417.16			Mean	0	0	-380.1
N2C1W1_K	55	Best	-1	0	-349	N2C1W4_A	73	Best	0	0	-290
		SD	0.29	0	20.9			SD	0	0	18.02
		Mean	-1	0	-450.29			Mean	0	0	-370.98
N2C1W1_L	55	Best	0	0	-278	N2C1W4_B	71	Best	0	0	-285
		SD	0	0	20.54			SD	0	0	13.56
		Mean	0	0	-390.78			Mean	0	0	-367.66
N2C1W1_M	48	Best	-15	0	-313	N2C1W4_C	77	Best	-1	0	-285
		SD	0.62	0	14.45			SD	1.10	0	13.56
		Mean	-17	0	-378.27			Mean	-2	0	-367.66
N2C1W1_N	48	Best	0	0	-289	N2C1W4_D	82	Best	0	0	-233
		SD	0	0	15.36			SD	0	0	18.25
		Mean	0	0	-366.61			Mean	0	0	-342.56
N2C1W1_O	54	Best	0	0	-268	N2C1W4_E	83	Best	0	0	-349
		SD	0	0	15.42			SD	0	0	20.9
		Mean	0	0	-366			Mean	0	0	-450.29
N2C1W1_P	54	Best	0	0	-293	N2C1W4_F	77	Best	-4	0	-278
		SD	0	0	13.26			SD	0.65	0	20.54
		Mean	0	0	-360.45			Mean	-4	0	-390.78
N2C1W1_Q	46	Best	-22	0	-315	N2C1W4_G	71	Best	-1	0	-326
		SD	2.03	0	14.14			SD	0.67	0	14.256
		Mean	-28	0	-379.81			Mean	-1	0	-398.36
N2C1W1_R	45	Best	-1	0	-233	N2C1W4_H	75	Best	0	0	-268
		SD	0.20	0	18.19			SD	0.77	0	15.25
		Mean	-1	0	-342.45			Mean	-1	0	-312
N2C1W1_S	45	Best	0	0	-349	N2C1W4_I	73	Best	0	0	-289
		SD	0	0	20.9			SD	0	0	12.36
		Mean	0	0	-450.29			Mean	0	0	-378
N2C1W1_T	52	Best	0	0	-313	N2C1W4_J	74	Best	0	0	-313
		SD	0	0	14.45			SD	0	0	14.256
		Mean	0	0	-378.27			Mean	0	0	-369.36

Table 3
Performance of proposed HS variations for 60 datasets with 100 objects (continued)

Dataset	Minimal bins	HS Variation			Dataset	Minimal bins	HS Variation				
		V3	V2	V1			V3	V2	V1		
N2C1W2_A	64	Best	0	0	-278	N2C1W4_K	70	Best	-11	0	-349
		SD	0	0	20.36			SD	3.80	0	20.9
		Mean	0	0	-398.78			Mean	-19	0	-450.29
N2C1W2_B	61	Best	-1	0	-369	N2C1W4_L	75	Best	0	0	-323
		SD	1.37	0	14.85			SD	0	0	14.25
		Mean	-3	0	-426.36			Mean	0	0	-378.26
N2C1W2_C	68	Best	0	0	-328	N2C1W4_M	72	Best	0	0	-307
		SD	0	0	15.46			SD	0	0	21.25
		Mean	0	0	-384.56			Mean	0	0	-416.36
N2C1W2_D	74	Best	0	0	-277	N2C1W4_N	71	Best	0	0	-298
		SD	0	0	20.98			SD	0	0	13.25
		Mean	0	0	-374.38			Mean	0	0	-369.36
N2C1W2_E	65	Best	-12	0	-532	N2C1W4_O	80	Best	0	0	-289
		SD	3.22	0	21.06			SD	0	0	13.366
		Mean	-22	0	-609.36			Mean	0	0	-369.36
N2C1W2_F	65	Best	0	0	-308	N2C1W4_P	67	Best	0	0	-295
		SD	0	0	21.72			SD	0	0	13.56
		Mean	0	0	-417.16			Mean	0	0	-369.66
N2C1W2_G	73	Best	0	0	-349	N2C1W4_Q	75	Best	-1	0	-287
		SD	0	0	20.9			SD	0	0	20.54
		Mean	0	0	-450.29			Mean	-1	0	-396.36
N2C1W2_H	70	Best	0	0	-289	N2C1W4_R	70	Best	0	0	-268
		SD	0	0	15.36			SD	2.36	0	15.25
		Mean	0	0	-366.61			Mean	-2	0	-369
N2C1W2_I	67	Best	-2	0	-313	N2C1W4_S	80	Best	0	0	-308
		SD	0.07	0	14.45			SD	0	0	21.72
		Mean	-2	0	-378.26			Mean	0	0	-417.26
N2C1W2_J	67	Best	0	0	-313	N2C1W4_T	70	Best	0	0	-349
		SD	0	0	28.46			SD	0	0	20.9
		Mean	0	0	-379.26			Mean	0	0	-450.29

Results from table 4 show that the performance of all three algorithms reduced with an increase in problem complexity with rise in problem size. Best, mean and standard deviations of all three algorithms reduced for instances with 200 objects. For few instances the performance of variation 2 and 3 are similar. There are many instances where variation 2 is found better than variation 3.

Table 4
Best, mean and standard deviation values of proposed HS variations for 60 datasets with 200 objects

Dataset	Minimal bins	HS Variation			Dataset	Minimal bins	HS Variation				
		V3	V2	V1			V3	V2	V1		
N3C1W1_A	105	Best	-34	-2	-1374	N3C1W2_K	120	Best	-43	-14	-1392
		SD	17.58	49.84	22.58			SD	18.30	43.37	32.26
		Mean	-84	-41	-1353			Mean	-93	-41	-1494.85
N3C1W1_B	114	Best	-9	0	-1523	N3C1W2_L	136	Best	-4	0	-1523
		SD	20.66	46.92	35.79			SD	18.78	45.69	35.79
		Mean	-65	-12.49	-1667.69			Mean	-49	-12.59	-1667.69
N3C1W1_C	99	Best	0	0	-1213	N3C1W2_M	136	Best	0	0	-1404
		SD	0	56.71	30.93			SD	22.01	46.06	34.4
		Mean	0	-17.51	-1354.46			Mean	-48	-17.69	-1567.06
N3C1W1_D	108	Best	-3	0	-1404	N3C1W2_N	136	Best	-18	0	-1428
		SD	16.14	57.83	34.4			SD	30.03	46.92	26.4
		Mean	-47	-19.04	-1567.06			Mean	-105	-19.04	-1534.86
N3C1W1_E	98	Best	-3	0	-1297	N3C1W2_O	127	Best	-6	0	-1374
		SD	6.46	70.4	29.74			SD	17.89	49.84	28.46
		Mean	-15	-28.75	-1392.08			Mean	-52	-28.75	-1471.27
N3C1W1_F	113	Best	-2	0	-1443	N3C1W2_P	126	Best	-1	0	-1235
		SD	14.03	66.12	28.16			SD	20.03	54.12	32.63
		Mean	-22	-21.84	-1535.08			Mean	-45	-21.65	-1397.12
N3C1W1_G	111	Best	0	0	-1392	N3C1W2_Q	135	Best	0	0	-1320
		SD	0	43.37	32.26			SD	23.08	56.71	32.35
		Mean	0	-10.76	-1494.85			Mean	-48	-10.25	-1448.88
N3C1W1_H	104	Best	-10	0	-1416	N3C1W2_R	123	Best	-52	-8	-1689
		SD	6.67	45.69	33.33			SD	22.14	57.83	37.44
		Mean	-26	-12.32	-1530.17			Mean	-120	-12.32	-1867.86
N3C1W1_I	100	Best	-1	0	-1310	N3C1W2_S	130	Best	0	0	-1404
		SD	14.52	46.05	29.35			SD	12.49	66.12	34.4
		Mean	-20	-10.47	-1411.53			Mean	-20	-10.48	-1567.06
N3C1W1_J	108	Best	0	0	-1374	N3C1W2_T	136	Best	-1	0	-1218
		SD	15.24	55.172	28.46			SD	14.12	72.4	30.68
		Mean	-21	-15.6	-1471.27			Mean	-22	-15.2	-1256.23

Table 4

Best, mean and standard deviation values of proposed HS variations for 60 datasets with 200 objects (Continued)

Dataset	Minimal bins	HS Variation			Dataset	Minimal bins	HS Variation				
		V3	V2	V1			V3	V2	V1		
N3C1W1_K	102	Best	-4	0	-1428	N3C1W4_A	149	Best	-8	-2	-1374
		SD	11.35	52.65	26.4			SD	17.56	43.37	22.58
		Mean	-41	-15.43	-1534.86			Mean	-51	-41	-1455.52
N3C1W1_L	97	Best	-50	0	-1235	N3C1W4_B	149	Best	-1	0	-1423
		SD	6.91	43.37	32.63			SD	14.11	45.69	33.58
		Mean	-69	-10.76	-1397.12			Mean	-22	-12.49	-1557.59
N3C1W1_M	106	Best	0	0	-1320	N3C1W4_C	146	Best	-12	0	-1425
		SD	14.38	45.69	32.35			SD	23.17	46.05	32.22
		Mean	-21	-12.32	-1448.88			Mean	-73	-17.51	-1597.19
N3C1W1_N	93	Best	-38	0	-1191	N3C1W4_D	148	Best	-13	0	-1566
		SD	8.35	46.05	28.35			SD	20.16	46.92	28.54
		Mean	-53	-10.47	-1311.8			Mean	-65	-19.04	-1452.33
N3C1W1_O	98	Best	-79	-12	-1311	N3C1W4_E	142	Best	-2	0	-1198
		SD	12.77	52.65	28.54			SD	11.42	49.84	25.36
		Mean	-117	-15.43	-1452.22			Mean	-25	-28.75	-1422.6
N3C1W1_P	108	Best	-10	0	-1432	N3C1W4_F	140	Best	-8	0	-1297
		SD	20.53	55.172	36.22			SD	14.68	55.172	29.74
		Mean	-42	-15.6	-1594.19			Mean	-44	-21.84	-1392.08
N3C1W1_Q	100	Best	-13	0	-1404	N3C1W4_G	148	Best	-10	0	-1215
		SD	6.23	45.69	34.4			SD	19.16	56.71	30.93
		Mean	-36	-12.32	-1567.06			Mean	-69	-10.76	-1354.26
N3C1W1_R	99	Best	-2	0	-1286	N3C1W4_H	141	Best	-4	0	-1678
		SD	12.67	0	28.26			SD	16.53	57.83	36.8
		Mean	-31	0	-1422.7			Mean	-62	-12.32	-1886.2
N3C1W1_S	100	Best	-27	0	-1269	N3C1W4_I	140	Best	-24	0	-1687
		SD	15.52	0	31.38			SD	20.29	66.12	37.55
		Mean	-69	0	-1425.03			Mean	-86	-10.47	-1879.36
N3C1W1_T	102	Best	-15	0	-1361	N3C1W4_J	142	Best	-21	0	-1523
		SD	13.62	0	29.24			SD	20.81	70.4	35.79
		Mean	-62	0	-1496.87			Mean	-76	-15.6	-1667.59
N3C1W2_A	125	Best	-2	-111	-1523	N3C1W4_K	147	Best	0	0	-1143
		SD	11.59	169.469	35.79			SD	0	66.12	28.25
		Mean	-28	-157.02	-1667.69			Mean	0	-80.58	-1535.09
N3C1W2_B	126	Best	0	-113	-1374	N3C1W4_L	148	Best	0	0	-1416
		SD	12.81	196.42	22.58			SD	22.48	56.71	33.33
		Mean	-19	-182.3	-1466.59			Mean	-45	-10.76	-1530.18
N3C1W2_C	125	Best	-42	-98	-1679	N3C1W4_M	149	Best	-8	0	-1404
		SD	14.47	162.689	26.7			SD	29.17	57.83	34.4
		Mean	-86	-140.53	-1886.22			Mean	-98	-12.32	-1567.06
N3C1W2_D	139	Best	0	-111	-1311	N3C1W4_N	148	Best	-4	0	-1392
		SD	15.81	169.469	28.54			SD	24.88	45.69	32.26
		Mean	-22	-157.02	-1452.22			Mean	-61	-12.49	-1494.85
N3C1W2_E	132	Best	-1	-1	-1191	N3C1W4_O	143	Best	-1	0	-1374
		SD	18.62	163.369	28.35			SD	15.78	70.4	22.58
		Mean	-47	-74.6	-1311.8			Mean	-34	-15.6	-1465.69
N3C1W2_F	123	Best	-30	-37	-1297	N3C1W4_P	143	Best	-3	0	-1298
		SD	19.30	217.29	29.74			SD	22.92	46.05	29.85
		Mean	-84	-106.14	-1392.08			Mean	-46	-17.51	-1326.35
N3C1W2_G	132	Best	-14	-30	-1213	N3C1W4_Q	146	Best	-23	0	-1215
		SD	13.85	174.79	30.93			SD	18.96	46.92	30.36
		Mean	-72	-87.75	-1354.46			Mean	-78	-19.04	-1365.26
N3C1W2_H	129	Best	-39	-79	-1416	N3C1W4_R	145	Best	0	0	-1326.35
		SD	16.14	165.43	33.33			SD	18.04	55.172	32.36
		Mean	-86	-125.78	-1530.17			Mean	-22	-21.84	-1445.36
N3C1W2_I	126	Best	-25	-48	-1443	N3C1W4_S	145	Best	-21	0	-1452
		SD	13.19	164.277	28.16			SD	23.75	49.84	36.3
		Mean	-72	-132.43	-1535.08			Mean	-75	-28.75	-1598.36
N3C1W2_J	126	Best	-20	-1	-1374	N3C1W4_T	146	Best	-1	0	-1258
		SD	14.61	192.38	22.58			SD	12.36	43.37	30.36
		Mean	-59	-81.66	-1466.59			Mean	-21	-41	-1258.3

Fig. 8 shows the performance of the basic harmony search algorithm for sample datasets with 100 and 200 objects. Results show reduced performance with an increase in number of objects in bin packing problem. The basic HS algorithm is subjected to premature convergence leading to local optimal solution. Fig. 9 shows fast convergence of version 2 algorithm for sample datasets with 100 objects and 200 objects.

The performance of all three algorithms is tested by varying the HMCR values. The HMCR values are varied from 0.1 to 0.9. Fig. 10, 11(a) and 11(b) present the performance of version 1, 2 and 3 for changing HMCR values respectively. Version 1 algorithm shows better convergence with higher values of HMCR. Performance increases with HMCR values. Best performance is achieved with HMCR values 0.9 and 0.8. The basic harmony search (V1) algorithm fails to give optimal solution for bin packing problem instances.

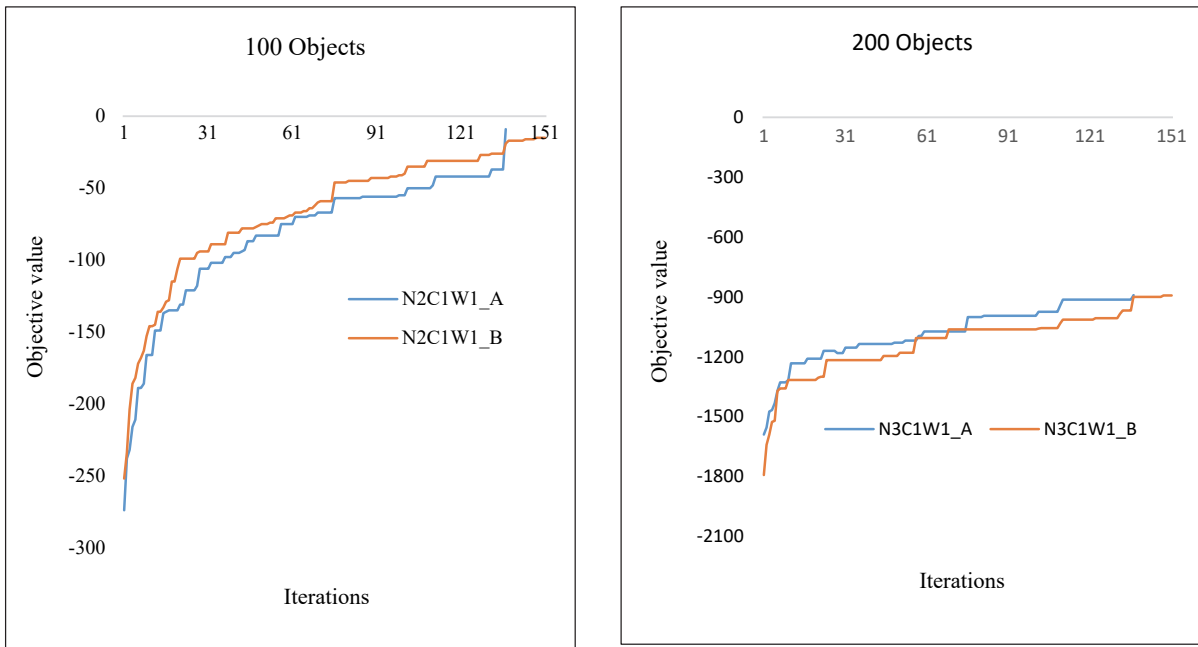


Fig. 8. Performance of version 1 algorithm for sample datasets with 100 and 200 objects

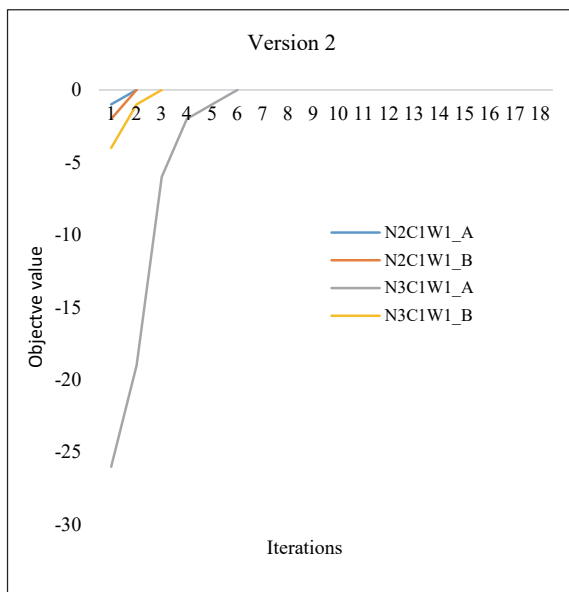


Fig. 9. Performance of version 2 algorithm for sample datasets with 100 and 200 objects

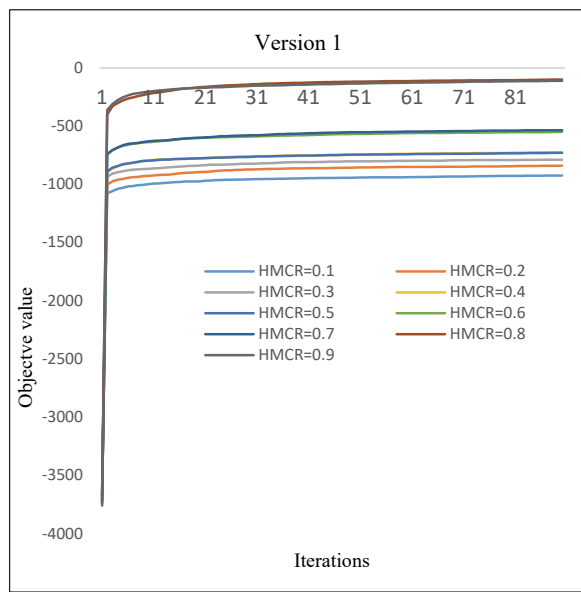
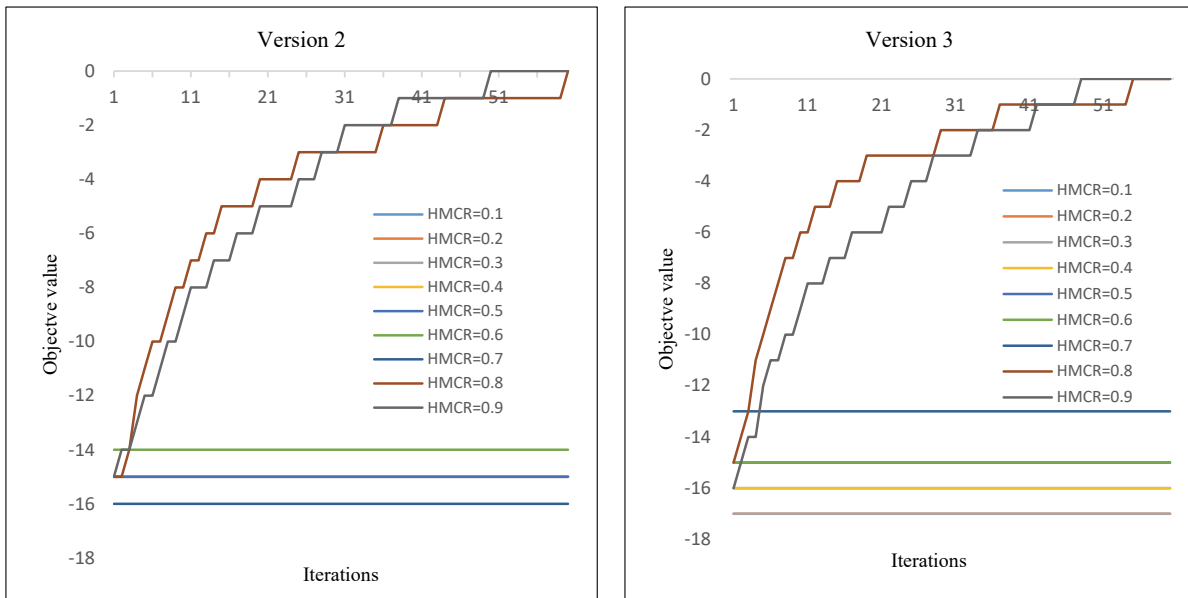


Fig. 10. Performance of version 1 algorithm for sample datasets with 100 and 200 objects

The performance of version 2 and version 3 is approximately same with changing HMCR values. Version 2 and 3 shows better convergence towards optimal solution. HMCR values 0.8 and 0.9 shows better performance as compared to other values. Harmony search algorithm with adaptive PAR shows better performance subjected to HMCR values.

Fig. 12 shows the effect of harmony memory size on performance of HS with adaptive PAR applied to solving bin packing problem. Impact of harmony memory size is very small on performance of algorithm.



(a) (b)
Fig. 11. Performance of algorithm for sample datasets with 100 and 200 objects (a) version 2 (b) version 3

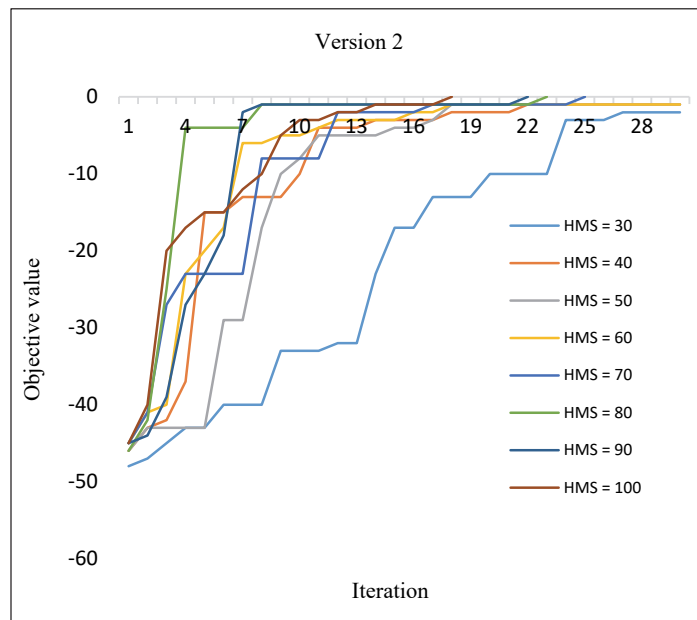


Fig. 12. Performance of version 2 algorithm for sample datasets with 100 and 200 objects

4. Conclusions

Bin packing problem is an important constraint satisfaction problem with huge combinations. In literature, many deterministic algorithms and heuristics algorithms are investigated to optimize the solutions. This paper presents the application of harmony search algorithm for solving bin packing problem.

Harmony search is a powerful heuristic to solve combinatorial optimization problem. It shows good convergence towards an optimal solution. Experiments show that HS stuck into the best region failing to reach optimal solution. To improve the performance of HS, paper presents two problem specific initialization strategies and re-initialization strategy. Results show that algorithm specific parameters have impact on the performance of HS to reach to optimal solution. Paper investigates two adaptive PAR strategies for fine tuning algorithm to optimize the solutions. The increasing PAR strategy is taken from (Mahdavi et al., 2007). Paper presents the results of basic and two variations of HS algorithm for solving 120 instances of

bin packing problem. Results show that the fine tuning of HS algorithm with adaptive PAR (increasing and decreasing) works better than static PAR. Performance of version 2 (increasing PAR) and version 3 (decreasing PAR) are similar with respect to best, mean and success rate for 60 bin packing problem instances with 100 objects. Version 2 is found better for larger problem instances with 200 objects. The performance of both HS variations are reduced with increase in problem complexity with size. Further experiments are conducted to analyze the impact of HMCR and HMS values of performance of algorithms. High HMCR values show better results. There is scope to improve the strategy for adaptive PAR for solving large instances on bin packing problem with minimal impact of HMCR values.

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